

Symbolic determinant identity testing and non-commutative ranks of matrix Lie algebras

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Abstract

One approach to make progress on the symbolic determinant identity testing (SDIT) problem is to study the structure of singular matrix spaces. After settling the non-commutative rank problem (Garg–Gurvits–Oliveira–Wigderson, *Found. Comput. Math.* 2020; Ivanyos–Qiao–Subrahmanyam, *Comput. Complex.* 2018), a natural next step is to understand singular matrix spaces whose non-commutative rank is full. At present, examples of such matrix spaces are mostly sporadic, so it is desirable to discover them in a more systematic way.

In this paper, we make a step towards this direction, by studying the family of matrix spaces that are closed under the commutator operation, that is, matrix Lie algebras. On the one hand, we demonstrate that matrix Lie algebras over the complex number field give rise to singular matrix spaces with full non-commutative ranks. On the other hand, we show that SDIT of such spaces can be decided in deterministic polynomial time. Moreover, we give a characterization for the matrix Lie algebras to yield a matrix space possessing singularity certificates as studied by Lovász (*B. Braz. Math. Soc.*, 1989) and Raz and Wigderson (*Building Bridges II*, 2019).

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1 Introduction

1.1 Background and motivations

Matrix spaces.

Let \mathbb{F} be a field. We use $M(\ell \times n, \mathbb{F})$ to denote the linear space of $\ell \times n$ matrices over \mathbb{F} , and let $M(n, \mathbb{F}) := M(n \times n, \mathbb{F})$. The general linear group of degree n over \mathbb{F} is denoted by $GL(n, \mathbb{F})$. A subspace \mathcal{B} of $M(\ell \times n, \mathbb{F})$ is called a *matrix space*, denoted by $\mathcal{B} \leq M(\ell \times n, \mathbb{F})$. Given $B_1, \dots, B_m \in M(n, \mathbb{F})$, $\langle B_1, \dots, B_m \rangle$ is the linear span of the B_i 's. In algorithms, $\mathcal{B} \leq M(n, \mathbb{F})$ is naturally represented by a linear basis $B_1, \dots, B_m \in M(n, \mathbb{F})$.

Two major algorithmic problems about matrix spaces are as follows.



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43 **The symbolic determinant identity testing problem.**

44 For $\mathcal{B} \leq M(n, \mathbb{F})$, let $\text{mrk}(\mathcal{B})$ be the maximum rank over all matrices in \mathcal{B} . We say that \mathcal{B} is
 45 *singular*, if $\text{mrk}(\mathcal{B}) < n$. To decide whether \mathcal{B} is singular is known as the *symbolic determinant*
 46 *identity testing* (SDIT) problem. The *maximum rank problem* for \mathcal{B} then asks to compute
 47 $\text{mrk}(\mathcal{B})$. The complexity of SDIT depends on the underlying field \mathbb{F} . When $|\mathbb{F}| = O(1)$, SDIT
 48 is coNP-complete [3]. When $|\mathbb{F}| = \Omega(n)$, by the polynomial identity testing lemma [23, 25],
 49 SDIT admits a randomized efficient algorithm. To present a deterministic polynomial-time
 50 algorithm for SDIT is a major open problem in computational complexity, as that would
 51 imply strong circuit lower bounds by the seminal work of Kabanets and Impagliazzo[18].

52 **The shrunk subspace problem.**

53 For $\mathcal{B} \leq M(n, \mathbb{F})$ and $U \leq \mathbb{F}^n$, the image of U under \mathcal{B} is $\mathcal{B}(U) := \{Bu : B \in \mathcal{B}, u \in U\}$.
 54 We say that U is a *shrunk subspace* of \mathcal{B} , if $\dim(U) > \dim(\mathcal{B}(U))$. The problem of deciding
 55 whether \mathcal{B} admits a shrunk subspace is the *shrunk subspace problem*. The *non-commutative*
 56 *rank problem*¹ [11, 16] asks to compute $\text{ncrk}(\mathcal{B}) := \max\{\dim(U) - \dim(\mathcal{B}(U)) : U \leq \mathbb{F}^n\}$.
 57 That is, \mathcal{B} admits a shrunk subspace if and only if its non-commutative rank is not full,
 58 i.e. $< n$. This problem is known for its connections to invariant theory, linear algebra,
 59 graph theory, and quantum information. Major progress in the past few years lead to
 60 deterministic efficient algorithms for the shrunk subspace problem, one by Garg, Gurvits,
 61 Oliveira, and Wigderson over fields of characteristic 0 [11], and the other by Ivanyos, Qiao,
 62 and Subrahmanyam over any field [15, 16].

63 **Motivations of our investigation.**

64 Note that if a matrix space admits a shrunk subspace, then it has to be singular. However,
 65 there exist singular matrix spaces without shrunk subspaces. After settling the shrunk
 66 subspace problem [11, 16], such matrix spaces form a bottleneck for further progress on
 67 SDIT. Moreover, ideas from these works are not expected to directly generalize as it was
 68 recently shown that the space of singular matrices cannot be seen as the null-cone of any
 69 reductive group action [21]

70 Two classical examples of such subspaces are as follows [20].

71 ► **Example 1.** 1. Let $\Lambda(n, \mathbb{F})$ be the linear space of alternating matrices, namely matrices
 72 satisfying $\forall v \in \mathbb{F}^n, v^t A v = 0$.² When n is odd, $\Lambda(n, \mathbb{F})$ is singular, as every alternating
 73 matrix is of even rank. Furthermore, it is easy to verify that $\Lambda(n, \mathbb{F})$ does not admit
 74 shrunk subspaces.

2. Let $C_1, \dots, C_n \in \Lambda(n, \mathbb{F})$, and let $\mathcal{C} \leq M(n, \mathbb{F})$ consist of all the matrices of the form
 $[C_1 v, C_2 v, \dots, C_n v]$, over $v \in \mathbb{F}^n$. As C_i 's are alternating, we have

$$v^t [C_1 v, C_2 v, \dots, C_n v] = [v^t C_1 v, v^t C_2 v, \dots, v^t C_n v] = 0,$$

75 so \mathcal{C} is singular. In [10], it is shown that when $n = 4$, certain choices of C_i ensure that \mathcal{C}
 76 does not have shrunk subspaces.

¹ The name “non-commutative” rank comes from a natural connection between matrix spaces and symbolic matrices over skew fields; see [11, 16] for details.

² When \mathbb{F} is of characteristic not 2, a matrix is alternating if and only if it is skew-symmetric.

77 While there are further examples in [1, 6], the above two examples (and their certain
78 subspaces) have been studied most in theoretical computer science and combinatorics, such
79 as by Lovász [20] and Raz and Wigderson [22], due to their connections to matroids and
80 graph rigidity.

81 As far as we see from the above, examples of singular matrix spaces without shrunk
82 subspaces in the literature are sporadic. Therefore, it is desirable to discover more singular
83 matrix spaces without shrunk subspaces, hopefully in a more systematic way. This is the
84 main motivation of this present article.

85 Overview of our main results.

86 Noting that the linear space of skew-symmetric matrices is closed under the commutator
87 bracket, we set out to study matrix Lie algebras. Our main results can be summarized as
88 follows.

- 89 ■ First, we show that matrix Lie algebras over \mathbb{C} gives rise to a family of singular matrix
90 spaces without shrunk subspaces. This result, partly inspired by [9], vastly generalizes
91 the linear spaces of skew-symmetric matrices.
- 92 ■ Second, we present a deterministic polynomial-time algorithm to solve SDIT for matrix Lie
93 algebras over \mathbb{C} . This algorithm heavily relies on the structural theory of, and algorithms
94 for, Lie algebras.
- 95 ■ Third, we examine when matrix Lie algebras are of the form in Example 1 (2) as above,
96 giving representation-theoretic criteria for such matrix Lie algebras.

97 In the rest of this introduction, we detail our results.

98 1.2 Our results

99 Recall that $\mathcal{B} \leq M(n, \mathbb{F})$ is a *matrix Lie algebra*, if \mathcal{B} is closed under the commutator bracket,
100 i.e. for any $A, B \in \mathcal{B}$, $[A, B] := AB - BA \in \mathcal{B}$.

101 We have striven to make this introduction as self-contained as possible. In an effort to
102 make this article accessible to wider audience, we summarize notions and results on Lie
103 algebras and representations relevant to this paper in Appendices A, B, and C.

104 Two results and a message.

105 We first study shrunk subspaces of matrix Lie algebras over \mathbb{C} . To state our results, we need
106 the following notions.

107 Given a matrix space $\mathcal{B} \leq M(n, \mathbb{F})$, $U \leq \mathbb{F}^n$ is an *invariant subspace* of \mathcal{B} , if for any
108 $B \in \mathcal{B}$, $B(U) \subseteq U$. We say that \mathcal{B} is *irreducible*, if the only invariant subspaces of \mathcal{B} are
109 0 and \mathbb{F}^n . The above notions naturally apply to matrix Lie algebras. The matrix space
110 $\mathcal{B} = 0 \leq M(1, \mathbb{F})$ is called the *trivial irreducible* matrix Lie algebra.

111 In general, let $\mathcal{B} \leq M(n, \mathbb{F})$ be a matrix Lie algebra. Then there exists $A \in GL(n, \mathbb{F})$,
112 such that $A^{-1}\mathcal{B}A$ is of block upper-triangular form, and each block on the diagonal defines
113 an irreducible matrix Lie algebra, called a *composition factor* of \mathcal{B} . Such an A defines a
114 chain of subspaces, called a *composition series* of the matrix Lie algebra \mathcal{B} . By the Jordan-
115 Hölder theorem, the isomorphic types of the composition factors are the same for different
116 composition series.

117 We then have the following criteria for the existence of shrunk subspaces of matrix Lie
118 algebras over \mathbb{C} .

119 ► **Theorem 2.** *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a non-trivial irreducible matrix Lie algebra. Then \mathcal{B} does*
 120 *not have a shrunk subspace.*

121 *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a matrix Lie algebra. Then \mathcal{B} has a shrunk subspace, if and only if*
 122 *one of its composition factors is the trivial matrix Lie algebra.*

123 The proof of Theorem 2 for the irreducible case makes use of the connection of Lie algebras
 124 and Lie groups as summarized in Appendix B. Going from the irreducible to the general case,
 125 we prove some basic properties of shrunk subspaces which may be of independent interest in
 126 Section 2.

127 After we proved Theorem 2, we learnt that Derksen and Makam independently proved it
 128 using a different approach via representation theory of Lie algebras [7].

129 We then present a deterministic polynomial-time algorithm to solve SDIT for matrix Lie
 130 algebras over \mathbb{C} . Our model of computation over \mathbb{C} will be explained in Section 4.

131 ► **Theorem 3.** *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a matrix Lie algebra. Then there is a deterministic*
 132 *polynomial-time algorithm to solve the symbolic determinant identity testing problem for \mathcal{B} .*

133 We believe that the strategy for the algorithm in Theorem 3 is interesting. It rests on the
 134 key observation that the maximum rank of \mathcal{B} is equal to the maximum rank of a Cartan
 135 subalgebra of \mathcal{B} . (We collect the notions and results on Cartan algebras relevant to this
 136 paper in Appendix C.) We then resort to the algorithm computing a Cartan subalgebra by de
 137 Graaf, Ivanyos and Rónyai [5] to get one. As Cartan subalgebras are upper-triangularisable,
 138 an SDIT algorithm can be devised easily.

139 Theorems 2 and 3 together bring out *the main message in this paper*: we identify non-
 140 trivial irreducible matrix Lie algebras over \mathbb{C} as an interesting families of matrix spaces,
 141 as (1) they do not admit shrunk subspaces, and (2) SDIT for such spaces can be solved in
 142 deterministic polynomial time.

143 To see that matrix Lie algebras do form an interesting family for the maximum rank
 144 problem, we list some examples.

145 ► **Example 4.** 1. Note that $\Lambda(n, \mathbb{F})$ is closed under the commutator bracket. Indeed, $\Lambda(n, \mathbb{F})$
 146 together with the commutator bracket is well-known as the orthogonal Lie algebra, and
 147 it is easy to see that it is irreducible.

148 2. Representations of abstract Lie algebras give rise to matrix Lie algebras. For example,
 149 let $\mathfrak{sl}(n, \mathbb{C})$ be the special linear Lie algebra, i.e, the Lie algebra of all $n \times n$ complex
 150 matrices with trace 0. Let $E_{i,j}$ be the elementary matrix with the only non-zero entry
 151 being 1 in the $(i, j)^{th}$ entry. A linear basis of $\mathfrak{sl}(n, \mathbb{C})$ consists of $E_{i,j}$, $i \neq j$. Consider
 152 for any fixed d , the vector space V spanned by all degree dn monomials in the variables
 153 $\{x_1, \dots, x_n\}$. Then, the representation is defined as $\rho(E_{ij})(x_1^{e_1} \dots x_n^{e_n}) = x_i \frac{\partial(x_1^{e_1} \dots x_n^{e_n})}{\partial x_j}$.

154 This gives rise to an irreducible matrix Lie algebra in $M(\binom{dn+n-1}{n-1}, \mathbb{C})$.

155 One may wonder whether irreducible matrix Lie algebras encompass singular and non-
 156 singular matrix spaces. To see this, note that $\Lambda(n, \mathbb{F})$ (as defined in Example 4) can be
 157 singular (for odd n) or non-singular (for even n). In fact, there is a representation-theoretic
 158 explanation for the maximum rank of certain irreducible matrix Lie algebras via weight
 159 spaces (Fact 34) as already observed by Draisma [8], from which it is evident that irreducible
 160 matrix Lie algebras can be singular or non-singular.

161 **Other singularity witnesses and matrix Lie algebras.**

162 After Theorem 2 and 3, we study further properties of matrix Lie algebras related to
 163 singularity as follows. Let $\mathcal{B} = \langle B_1, \dots, B_m \rangle \leq M(n, \mathbb{F})$ be a matrix space. Let x_1, \dots, x_m

164 be a set of commutative variables. Then $B = x_1B_1 + \dots + x_mB_m$ is a matrix of linear forms
 165 in x_i 's. When \mathbb{F} is large enough, the singularity of \mathcal{B} is equivalent to that of B over the
 166 function field. Viewing B as a matrix over the rational function field $\mathbb{F}(x_1, \dots, x_n)$, its kernel
 167 is spanned by vectors whose entries are polynomials. Let $v \in \mathbb{F}[x_1, \dots, x_m]^n$ be in $\ker(B)$.
 168 By splitting v according to degrees if necessary, we can assume that v is *homogeneous*, i.e.
 169 each component of v is homogeneous of degree d .

170 We are interested in those vectors in the kernel whose entries are linear forms. This is
 171 also partly motivated by understanding witnesses for singularity of matrix spaces, as by
 172 [18], putting SDIT in $\text{NP} \cap \text{coNP}$ already implies strong circuit lower bounds. Suppose B
 173 admits $v \in \ker(B)$ whose components are homogeneous degree- d polynomials. Then, ignoring
 174 bit complexities, v is a singularity witness of \mathcal{B} of size $O(m^d \times n)$, and the existence of a
 175 certificate of degree d can be checked and found in time $O(m^d \times n)$, by writing a linear
 176 system in $O(m^d \times n)$ variables.

177 Let $v_1, \dots, v_m \in \mathbb{F}^n$, and $v = x_1v_1 + \dots + x_mv_m$ be a vector of linear forms. We say
 178 that v is a (left homogeneous) *linear kernel vector* of B , if each entry of v^tB is the zero
 179 polynomial. Similarly, v is a right homogeneous linear kernel vector, if each entry of Bv is
 180 the zero polynomial.

181 Clearly, whether such a nonzero v exists does not depend on the choice of bases. Indeed,
 182 we can give a basis-free definition of a linear kernel vector for a matrix space $\mathcal{B} \leq M(n, \mathbb{F})$ as
 183 a non-zero linear map $\psi : \mathcal{B} \rightarrow \mathbb{F}^n$ such that for each A , $\psi(A)^tA = 0$.

184 Matrix spaces with linear kernel vectors have appeared in papers by Lovász [20] and Raz
 185 and Wigderson [22]. To see this, note that matrix spaces with linear kernel vectors can be
 186 constructed from alternating matrices as exhibited in Example 1 (2).

187 One approach for Lie algebras to yield matrix spaces with linear kernel vectors is through
 188 adjoint representations.

189 Recall that, given a Lie algebra $[-, -] : \mathfrak{g} \times \mathfrak{g} \rightarrow \mathfrak{g}$, the adjoint representation of \mathfrak{g} is
 190 $\text{ad} : \mathfrak{g} \rightarrow \mathfrak{gl}(\mathfrak{g})$ defined as $\text{ad}_x(y) = [x, y]$ for $x, y \in \mathfrak{g}$. The image of ad is a matrix space
 191 $\mathcal{A} \leq M(d, \mathbb{F})$ where $d = \dim(\mathfrak{g})$. As the Lie bracket $[\cdot, \cdot]$ is alternating, \mathcal{A} admits a linear
 192 kernel vector by the construction in Example 1 (2).

193 Our next theorem characterizes Lie algebra representations with linear kernel vectors.
 194 (We collect some basic notions of Lie algebra representations relevant to this paper in
 195 Appendix A.) Since we are concerned with matrix spaces which are images of Lie algebra
 196 representations, i.e. $\mathcal{B} = \rho(\mathfrak{g})$ where ρ is a representation of the Lie algebra \mathfrak{g} , we can assume
 197 without generality that ρ is faithful.

198 ► **Theorem 5.** *Let \mathcal{B} be the image of a faithful irreducible representation ϕ of a semisimple*
 199 *Lie algebra \mathfrak{g} over algebraically closed fields of characteristic not 2 or 3. Then \mathcal{B} admits*
 200 *a linear kernel vector if and only if \mathcal{B} is trivial, or \mathfrak{g} is simple and ϕ is isomorphic to the*
 201 *adjoint representation.*

202 1.3 Open questions.

203 Several questions can be asked after this work. First, can we identify more families of singular
 204 matrix spaces without shrunk subspaces? Second, our algorithm for SDIT of matrix Lie
 205 algebras heavily relies on the structure theory of Lie algebras and works over \mathbb{C} . It will
 206 be interesting to devise an alternative algorithm that is of a different nature, and works
 207 for matrix Lie algebras over fields of positive characteristics. Third, characterize those
 208 representations of non-semisimple Lie algebras with linear kernel vectors.

209 **The structure of the paper.**

210 In Section 2 we prove some results on shrunk subspaces that will be useful to prove Theorem 2.
 211 In Section 3 we prove Theorem 2. In Section 4 we prove Theorem 3. In Section 5 we prove
 212 Theorem 5.

213 **2 On shrunk subspaces of matrix spaces**

214 In this section we present some basic results and properties regarding shrunk subspaces and
 215 non-commutative ranks of matrix spaces.

216 **2.1 Canonical shrunk subspaces**

217 Let $\mathcal{B} \leq M(n, \mathbb{F})$. For a subspace U of \mathbb{F}^n define $\text{sd}_{\mathcal{B}}(U)$ as the difference $\dim(U) - \dim(\mathcal{B}(U))$.
 218 Thus $\text{sd}_{\mathcal{B}}(U)$ is positive for a shrunk subspace U and negative if \mathcal{B} expands U . We then have
 219 the following.

220 ► **Lemma 6.** *The function $\text{sd}_{\mathcal{B}}$ is supermodular. More specifically, if U_1 and U_2 are two
 221 subspaces of \mathbb{F}^n , then,*

$$222 \quad \text{sd}_{\mathcal{B}}(U_1 \cap U_2) + \text{sd}_{\mathcal{B}}(\langle U_1 \cup U_2 \rangle) \geq \text{sd}_{\mathcal{B}}(U_1) + \text{sd}_{\mathcal{B}}(U_2). \quad (1)$$

223 **Proof.** By modularity of the dimension, we have

$$224 \quad \dim(U_1 \cap U_2) + \dim(\langle U_1 \cup U_2 \rangle) = \dim(U_1) + \dim(U_2)$$

225 and

$$226 \quad \dim(\mathcal{B}(U_1) \cap \mathcal{B}(U_2)) + \dim(\langle \mathcal{B}(U_1) \cup \mathcal{B}(U_2) \rangle) = \dim(\mathcal{B}(U_1)) + \dim(\mathcal{B}(U_2)).$$

227 The second equality, using also that $\mathcal{B}(\langle U_1 \cup U_2 \rangle) = \langle \mathcal{B}(U_1) \cup \mathcal{B}(U_2) \rangle$ and $\mathcal{B}(U_1 \cap U_2) \leq$
 228 $\mathcal{B}(U_1) \cap \mathcal{B}(U_2)$, gives,

$$229 \quad \dim(\mathcal{B}(U_1 \cap U_2)) + \dim(\mathcal{B}(\langle U_1 \cup U_2 \rangle)) \leq \dim(\mathcal{B}(U_1)) + \dim(\mathcal{B}(U_2)).$$

230 Subtracting the last inequality from the first equality gives (1). ◀

231 ► **Proposition 7.** *Let $\mathcal{B} \leq M(n, \mathbb{F})$. Suppose $\text{ncrk}(\mathcal{B}) = n - c$ for $c > 0$. Then there exists
 232 a unique subspace $U \leq \mathbb{F}^n$ of the smallest dimension satisfying $\dim(U) - \dim(\mathcal{B}(U)) = c$,
 233 and there exists a unique subspace $U' \leq \mathbb{F}^n$ of the largest dimension such that $\dim(U') -$
 234 $\dim(\mathcal{B}(U')) = c$.*

235 **Proof.** We use the supermodular function $\text{sd}_{\mathcal{B}}$ defined in Lemma 6. Let U_1 and U_2 be
 236 subspaces with $\text{sd}_{\mathcal{B}}(U_i) = c$. Then Lemma 6 gives $\text{sd}_{\mathcal{B}}(U_1 \cap U_2) + \text{sd}_{\mathcal{B}}(\langle U_1 \cup U_2 \rangle) \geq 2c$.
 237 On the other hand, by the definition of the noncommutative rank, $\text{sd}_{\mathcal{B}}(U_1 \cap U_2) \leq c$ and
 238 $\text{sd}_{\mathcal{B}}(\langle U_1 \cup U_2 \rangle) \leq c$. It follows that all the three inequalities are in fact equalities. Thus
 239 the intersection as well as the span of all the subspaces U with $\text{sd}_{\mathcal{B}}(U) = c$ also have this
 240 property. ◀

241 By Proposition 7, in the case $\text{ncrk}(\mathcal{B}) = n - c$ for $c > 0$, we shall refer to the subspace U
 242 of the smallest dimension satisfying $\dim(U) - \dim(\mathcal{B}(U)) = c$ as the (lower) *canonical shrunk*
 243 *subspace*. The algorithm from [15, 16] actually computes the canonical shrunk subspace.

244 A natural group action on matrix spaces is as follows. Let $G = \text{GL}(n, \mathbb{F}) \times \text{GL}(n, \mathbb{F})$.
 245 Then $(A, C) \in G$ sends $\mathcal{B} \leq M(n, \mathbb{F})$ to $ABC^{-1} = \{ABC^{-1} : B \in \mathcal{B}\}$. The stabilizer group
 246 of this action on \mathcal{B} is denoted as $\text{Stab}(\mathcal{B}) = \{(A, C) \in G : ABC^{-1} = \mathcal{B}\}$. We then have the
 247 following proposition.

248 ► **Proposition 8.** *Let $\mathcal{B} \leq M(n, \mathbb{F})$. Suppose $\text{ncrk}(\mathcal{B}) = n - c^3$ for $c > 0$. Then for*
 249 *$\forall (A, C) \in \text{Stab}(\mathcal{B})$, the canonical shrunk subspace U is invariant under C , i.e., $C(U) = U$.*

250 **Proof.** From the definition of $\text{Stab}(\mathcal{B})$, we have $ABC^{-1} = \mathcal{B}$ and thus, $A\mathcal{B} = \mathcal{B}C$. Consider
 251 the subspace $A(U)$. Then, $\mathcal{B}(C(U)) = (\mathcal{B}C)(U) = (A\mathcal{B})(U) = A(\mathcal{B}(U))$. Since $A, C \in$
 252 $\text{GL}(n, \mathbb{F})$, $\dim(C(U)) = \dim(U)$ and $\dim(A(\mathcal{B}(U))) = \dim(\mathcal{B}(U))$. It follows that $C(U)$ is
 253 also a c -shrunk subspace of the same dimension as U . We then conclude that $C(U) = U$ by
 254 Proposition 7. ◀

255 2.2 Shrunk subspaces of block upper-triangular matrix spaces

256 Consider the following situation. Suppose $\mathcal{B} \leq M(n, \mathbb{F})$ satisfies that any $B \in \mathcal{B}$ is in the
 257 block upper-triangular form, i.e.

$$258 \quad B = \begin{bmatrix} C_1 & D_{1,2} & \dots & D_{1,d} \\ 0 & C_2 & \dots & D_{2,d} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & C_d \end{bmatrix},$$

259 where C_i is of size $n_i \times n_i$. Let

$$260 \quad \mathcal{C}_i = \langle C_i \in M(n_i, \mathbb{F}) : C_i \text{ appears as the } i\text{th diagonal block of some } B \in \mathcal{B} \rangle.$$

261 ► **Lemma 9.** *Let $\mathcal{B} \leq M(n, \mathbb{F})$, and let $V \leq \mathbb{F}^n$ such that $\mathcal{B}(V) \leq V$. If there exists a shrunk*
 262 *subspace for \mathcal{B} , then there also exist one which is either included in V or contains V .*

263 **Proof.** Assume that V itself is not a shrunk subspace. Then $\mathcal{B}(V) = V$. Let U be a shrunk
 264 subspace of \mathcal{B} . By Lemma 6, we have $\text{sd}_{\mathcal{B}}(V \cap U) + \text{sd}_{\mathcal{B}}(\langle V \cup U \rangle) \geq \text{sd}_{\mathcal{B}}(V) + \text{sd}_{\mathcal{B}}(U) =$
 265 $0 + \text{sd}_{\mathcal{B}}(U) > 0$. Thus either $\text{sd}_{\mathcal{B}}(V \cap U)$ or $\text{sd}_{\mathcal{B}}(\langle V \cup U \rangle)$ must be positive. ◀

266 The following proposition characterizes the existence of shrunk subspaces in block upper-
 267 triangular matrix spaces.

268 ► **Proposition 10.** *Let $\mathcal{B} \leq M(n, \mathbb{F})$ and $\mathcal{C}_i \leq M(n_i, \mathbb{F})$, $i \in [d]$, as above. Then \mathcal{B} has a*
 269 *shrunk subspace if and only if there exists $i \in [d]$ such that \mathcal{C}_i has a shrunk subspace.*

270 **Proof.** The if direction can be verified easily. For the only if direction, we induct on d . When
 271 $d = 1$, this is clear. Suppose this holds for $d < k$. Consider $d = k$, and suppose \mathcal{B} admits
 272 a shrunk subspace. Let $V \leq \mathbb{F}^n$ be the subspace spanned by those standard basis vectors
 273 $e_{n_1+1}, e_{n_1+2}, \dots, e_n$. We then have two cases.

- 274 1. There exists a shrunk subspace $W \leq V$. In this case, by the induction hypothesis, there
 275 exists $i \in \{2, \dots, n\}$ such that \mathcal{C}_i has a shrunk subspace.
- 276 2. There are no shrunk subspaces $W \leq V$. Then by Lemma 9, there exists a shrunk subspace
 277 W such that $W > V$. Then by considering W/V , we obtain a shrunk subspace for \mathcal{C}_1 .

278 This concludes the proof of Proposition 10. ◀

³ Recall that $\text{ncrk}(\mathcal{B}) := \max\{\dim(U) - \dim(\mathcal{B}(U)) : U \leq \mathbb{F}^n\}$.

279 **3 Shrunken subspaces of matrix Lie algebras over \mathbb{C}**

280 In this section, we will give a characterization of those matrix Lie algebras over \mathbb{C} with
 281 shrunken subspaces, proving Theorem 2. The main reason for working over \mathbb{C} is to make use of
 282 the connections between Lie algebras and Lie groups as described in Appendix B.

283 We will first give such a characterization for irreducible matrix Lie algebras. The general
 284 case then follows by combining this with the results in Section 2.2.

285 The key to understanding the irreducible case lies in the following lemma; for notions
 286 such as matrix exponentiation and derivation, cf. Appendix B.

287 **► Lemma 11** ([12, Proposition 4.5 (1)]). *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be an irreducible matrix Lie algebra.
 288 Let $W \leq \mathbb{C}^n$ and $M \in \mathcal{B}$. If $e^{tM}(W) \leq W$ for all $t \in \mathbb{R}$, then $M(W) \leq W$.*

289 **Proof.** Take any $w \in W$. Note that $\frac{d(e^{tM})}{dt}(w) = (Me^{tM})(w) = M(e^{tM}(w))$, and $\frac{d(e^{tM})}{dt} =$
 290 $\lim_{h \rightarrow t} \frac{e^{hM} - e^{tM}}{h - t}$. So at $t = 0$, we have $M(w) = \lim_{h \rightarrow 0} \frac{e^{hM}(w) - w}{h}$. Since $e^{tM}(w) \in W$ for all
 291 $t \in \mathbb{R}$, $\frac{e^{hM}(w) - w}{h}$ lies in W for any h , and so does the limit which is $M(w)$. ◀

292 We will also need the following result.

293 **► Lemma 12.** *Given a matrix Lie algebra $\mathcal{B} \leq M(n, \mathbb{C})$, we have that $\forall t \in \mathbb{R}$ and $M \in \mathcal{B}$,
 294 $e^{tM}\mathcal{B}e^{-tM} = \mathcal{B}$.*

295 **Proof.** By the connection between Lie groups and Lie algebras (cf. Theorem 25), there
 296 exists some Lie group G whose associated Lie algebra is \mathcal{B} . This implies that for any $M \in \mathcal{B}$,
 297 $e^{tM} \in G$. Then by the fact that the conjugation of $g \in G$ stabilizes \mathcal{B} (cf. Theorem 26), we
 298 have $e^{tM}\mathcal{B}e^{-tM} = \mathcal{B}$. ◀

299 We are now ready to prove Theorem 2.

300 **► Theorem 2.** *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a non-trivial irreducible matrix Lie algebra. Then \mathcal{B} does
 301 not have a shrunken subspace.*

302 *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a matrix Lie algebra. Then \mathcal{B} has a shrunken subspace, if and only if
 303 one of its composition factors is the trivial matrix Lie algebra.*

304 **Proof.** We first handle the irreducible case.

305 For the sake of contradiction, suppose \mathcal{B} has a shrunken subspace. Then let $V = \mathbb{C}^n$, and
 306 let $U \leq V$ be the canonical shrunken subspace of \mathcal{B} . By Lemma 12, for any $M \in \mathcal{B}$, we have
 307 that $(e^{tM}, e^{tM}) \in \text{Stab}(\mathcal{B})$. By Proposition 8, U is invariant under e^{tM} . By Lemma 11, U
 308 is an invariant subspace of \mathcal{B} .

309 Since \mathcal{B} is irreducible as a matrix Lie algebra, the only invariant subspaces are 0 and V .
 310 Since U is a shrunken subspace, it cannot be 0. If $U = V$, then $\mathcal{B}(V)$ is a proper subspace of V .
 311 If $\mathcal{B}(V)$ is non-zero, then $\mathcal{B}(\mathcal{B}(V)) \leq \mathcal{B}(V)$. This implies that $\mathcal{B}(V)$ is a proper and non-zero
 312 invariant subspace of \mathcal{B} , which is impossible as \mathcal{B} is irreducible. It follows that $U = V$ and
 313 $\mathcal{B}(V) = 0$. In this case, V must be of dimension 1, as any non-zero proper subspace of V is
 314 an invariant subspace. It follows that \mathcal{B} has to be the trivial matrix Lie algebra. We then
 315 arrive at the desired contradiction.

316 The general case follows from the irreducible case as shown above, and Proposition 10. ◀

317 **4 SDIT for matrix Lie algebras over \mathbb{C}**

318 In this section, we present a deterministic polynomial-time algorithm that solves SDIT for
319 matrix Lie algebras over \mathbb{C} , proving Theorem 3.

320 The basic idea is to realize that \mathcal{B} is singular if and only if every Cartan subalgebra of \mathcal{B}
321 is singular. Furthermore, a Cartan subalgebra is nilpotent, so in particular it is solvable. It
322 follows, by Lie's theorem (Theorem 28), that a Cartan subalgebra of a matrix Lie algebra is
323 upper-triangularisable by the conjugation action. A key task here is to compute a Cartan
324 subalgebra of \mathcal{B} . This problem has been solved by de Graaf, Ivanyos, and Rónyai in [5].

325 **Computation model over \mathbb{C} .**

326 We adopt the following computation model over \mathbb{C} , in consistent with that in [5]. That is,
327 we assume the input matrices are over a number field \mathbb{E} . Therefore \mathbb{E} is a finite-dimensional
328 algebra over \mathbb{Q} . If $\dim_{\mathbb{Q}}(\mathbb{E}) = d$, then \mathbb{E} is the extension of \mathbb{F} by a single generating element
329 α , so \mathbb{E} can be represented by the minimal polynomial of α over \mathbb{F} , together with an isolating
330 rectangle for α in the case of \mathbb{C} .

331 **4.1 Cartan subalgebras.**

332 We collect notions and results on Cartan subalgebras useful to us in Appendix C. Here, we
333 recall the following. Let \mathfrak{g} be a Lie algebra. A subalgebra $\mathfrak{h} \subseteq \mathfrak{g}$ is a *Cartan subalgebra*, if it
334 is nilpotent and self-normalizing.

335 In [5], de Graaf, Ivanyos, and Rónyai studied the problem of computing Cartan subalgebras.
336 We state the following version of their main result in our context as follows. For a more
337 precise statement, see Theorem 30.

338 **► Theorem 13** ([5, Theorem 5.8]). *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a matrix Lie algebra. Then there*
339 *exists a deterministic polynomial-time algorithm that computes a linear basis of a Cartan*
340 *subalgebra \mathcal{A} of \mathcal{B} .*

341 **4.1.1 Maximum ranks of Cartan subalgebras.**

342 The key lemma that supports our algorithm is the following.

343 **► Lemma 14.** *Let $\mathcal{B} \leq M(n, \mathbb{C})$ be a matrix Lie algebra. Let $\mathcal{A} \leq \mathcal{B}$ be a Cartan subalgebra.*
344 *Then, $\text{mrk}(\mathcal{B}) = \text{mrk}(\mathcal{A})$.*

345 **Proof.** We shall utilise two results about Cartan subalgebras; for details see Appendix C.

346 First, let \mathfrak{g} be a Lie algebra over a large enough field. Then there exists a set of
347 generic⁴ elements $R \subseteq \mathfrak{g}$, such that for any $x \in R$, the Fitting null component of ad_x ,
348 $F_0(\text{ad}_x) = \{y \in \mathfrak{g} : \exists m > 0, \text{ad}_x^m(y) = 0\}$, is a Cartan subalgebra. For a precise statement,
349 see Theorem 29.

350 Second, let \mathcal{B} be a matrix Lie algebra over \mathbb{C} . Then for any two Cartan subalgebras \mathcal{A} ,
351 \mathcal{A}' of \mathcal{B} , they are conjugate, namely there exists $T \in \text{GL}(n, \mathbb{C})$ such that $T\mathcal{A}T^{-1} = \mathcal{A}'$. For
352 a precise statement, see Theorem 27.

353 By the first result, in particular by the fact that elements in R are generic, there exists a
354 matrix $C \in \mathcal{B}$ of rank $\text{mrk}(\mathcal{B})$, such that $\mathcal{C} := F_0(\text{ad}_C)$ is a Cartan subalgebra. Noting that

⁴ This means that after identifying \mathfrak{g} with $\mathbb{F}^{\dim(\mathfrak{g})}$, these elements form a Zariski open set.

355 $C \in \mathcal{C}$, $\text{mrk}(C) = \text{mrk}(\mathcal{B})$. By the second result, for any Cartan subalgebra \mathcal{A} of \mathcal{B} , \mathcal{A} and C
 356 are conjugate, which implies that $\text{mrk}(\mathcal{A}) = \text{mrk}(C) = \text{mrk}(\mathcal{B})$. ◀

357 **4.2 Upper-triangularisable matrix spaces.**

358 Let $\mathcal{B} \leq M(n, \mathbb{F})$. We say that \mathcal{B} is upper-triangularisable, if there exists $S, T \in GL(n, \mathbb{F})$,
 359 such that for any $B \in \mathcal{B}$, SBT is upper-triangular. Upper-triangularisable matrix spaces
 360 are of interest to us, because solvable matrix Lie algebras can be made simultaneously
 361 upper-triangular via conjugation by Lie’s theorem (Theorem 28).

362 If a matrix space is upper-triangularisable, then we can decide if \mathcal{B} is singular in a
 363 black-box fashion, as its singularity is completely determined by the diagonals of the resulting
 364 upper-triangular matrix space. The following lemma is well-known and we include a proof
 365 for completeness.

366 ▶ **Lemma 15.** *Let $n, k \in \mathbb{N}$. Let \mathbb{F} be a field such that $|\mathbb{F}| > (k - 1)n$. There exists a
 367 deterministic algorithm that outputs in time $\text{poly}(n, k)$ a set $\mathcal{H} \subseteq \mathbb{F}^k$, such that any non-zero
 368 k -variate degree- n polynomial, which is a product of linear forms, evaluates to a non-zero
 369 value on at least one point in \mathcal{H} .*

370 **Proof.** Let ℓ_1, \dots, ℓ_n be n non-zero linear forms in k variables. We can also identify them
 371 as vectors in \mathbb{F}^k by taking their coefficients. Fix a subset $S \subseteq \mathbb{F}$ of size $(k - 1)n + 1$. Let
 372 $\mathcal{H} = \{(1, \alpha, \dots, \alpha^{k-1}) \mid \alpha \in S\}$. This is clearly a set of size $(k - 1)n + 1$.

373 We claim that any non-zero linear form ℓ_i vanishes on at most $k - 1$ points in \mathcal{H} . This
 374 is because if it vanishes on k points, we have $A\ell_i = 0$ where A is the Vandermonde matrix
 375 corresponding to those k points. This is impossible because the Vandermonde matrix is
 376 invertible and ℓ_i is non-zero.

377 It follows that there is at least one point in \mathcal{H} such that every ℓ_i has a non-zero evaluation
 378 at this point. This concludes the proof. ◀

380 **4.3 The algorithm.**

381 Given the above preparation, we present the following algorithm for computing the commut-
 382 ative rank of a matrix Lie algebra.

383 **Input:** $\mathcal{B} = \langle B_1, \dots, B_m \rangle \leq M(n, \mathbb{C})$, such that \mathcal{B} is a matrix Lie algebra.

384 **Output:** “Singular” if \mathcal{B} is singular, and “Non-singular” otherwise.

385 **Algorithm: 1.** Use Theorem 13 to obtain $\mathcal{C} = \langle C_1, \dots, C_k \rangle \leq \mathcal{B}$, such that \mathcal{C} is a Cartan
 386 subalgebra of \mathcal{B} .

387 2. Use Lemma 15 to obtain $H \subseteq \mathbb{C}^k$, $|H| = (k - 1)n + 1$.

388 3. For any $(\alpha_1, \dots, \alpha_k) \in H$, if $\sum_{i \in [k]} \alpha_i C_i$ is non-singular, return “Non-singular”.

389 4. Return “Singular”.

390 The above algorithm clearly runs in polynomial time. The correctness of the above
 391 algorithm follows from Lemmas 14 and 15, as well as Lie’s theorem on solvable Lie
 392 algebras (Theorem 28). This concludes the proof of Theorem 3.

393 ▶ **Remark 16.** We do not solve the maximum rank problem for matrix Lie algebras in general.
 394 While the maximum rank problem for matrix Lie algebras reduces to the maximum rank
 395 problem for upper-triangularisable matrix spaces through Cartan subalgebras, to compute the
 396 maximum rank for the latter deterministically seems difficult. This is because the maximum

rank problem for upper-triangularisable matrix spaces is as difficult as the general SDIT problem, an observation already in [14].

There is one case where we do solve the maximum rank problem, that is, when the matrix Lie algebra over \mathbb{C} is semisimple. In this case, Cartan subalgebras are diagonalizable [13, Theorem in Chapter 6.4]. Therefore, in the above algorithm we can output the maximum rank over $\sum_{i \in [k]} \alpha_i C_i$ where $(\alpha_1, \dots, \alpha_k) \in H$ as the maximum rank of \mathcal{B} .

5 Linear kernel vectors of matrix Lie algebras

The goal of this section is to study existence of linear kernel vectors for matrix spaces arising from representations of Lie algebras.

Let \mathfrak{g} be a Lie algebra and (ρ, V) be a representation of \mathfrak{g} , where $V \cong \mathbb{F}^n$. Let $\mathcal{B} = \rho(\mathfrak{g}) \leq M(n, \mathbb{F})$.

First, note that \mathcal{B} admits a common kernel vector⁵ if and only if (ρ, V) has a trivial subrepresentation. We view this as a degenerate case, so in the following we shall mainly consider representations without trivial subrepresentations.

By the basis-free definition of linear kernel vectors in Section 1.2, $\mathcal{B} = \rho(\mathfrak{g})$ has a linear kernel vector if we have a linear map $\beta : \rho(\mathfrak{g}) \rightarrow V$ such that $\rho(x)\beta(\rho(x)) = 0$. For our purposes, it will be more convenient to work with a map from \mathfrak{g} itself to V . This leads us to define that for a representation (ρ, V) , the linear map $\psi : \mathfrak{g} \rightarrow V$ is a *linear kernel vector* if

$$\rho(x)\psi(x) = 0 \tag{2}$$

for every $x \in \mathfrak{g}$. We further assume that ψ is not identically zero.

► **Remark 17.** The definition of linear kernel vectors above is a generalization that allows for possibly more linear kernel vectors. This is because a linear kernel vector β yields a generalized one by taking $\psi = \beta \circ \rho$. However, when ρ is not injective, we can have many more generalized maps. For example, for a trivial representation $(0, V)$, β has to be 0 but any linear map from \mathfrak{g} to V is a generalized linear kernel vector.

Applying Equation (2) to x, y and $x + y$ one obtains for every $x, y \in \mathfrak{g}$,

$$\rho(x)\psi(y) + \rho(y)\psi(x) = 0 \tag{3}$$

Since $[x, x] = 0$ for the adjoint representation, the identity map of \mathfrak{g} and its scalar multiples are generalized linear kernel vectors.

Assume that $\psi : \mathfrak{g} \rightarrow V$ is a linear kernel vector for (ρ, V) . Let (ρ', V') be another representation of \mathfrak{g} . Then, if $\phi : V \rightarrow V'$ is a non-zero linear map such that $\phi \circ \rho = \rho' \circ \phi$ (that is, ϕ is a homomorphism between the two representations) then $\phi \circ \psi$ is a linear kernel vector for (ρ', V') . Indeed, $\rho'(x)\phi(\psi(x)) = \phi(\rho(x)(\psi(x))) = \phi(0) = 0$.

Our aim is to show that for many of Lie algebras \mathfrak{g} , unless the representation (ρ, V) includes a trivial subrepresentation, every linear kernel vector $\psi : \mathfrak{g} \rightarrow V$ can be obtained as the composition of the adjoint representation and a homomorphism.

► **Theorem 18.** *Let \mathfrak{g} be a semisimple Lie algebra \mathfrak{g} over an algebraically closed field \mathbb{F} of characteristic not 2 or 3. Assume that the trivial representation is not a subrepresentation of the representation (ρ, V) of \mathfrak{g} . Then any linear kernel vector ψ defines a homomorphism $\psi : (\text{ad}, \mathfrak{g}) \rightarrow (\rho, V)$ i.e., for every $x, y \in \mathfrak{g}$,*

$$\psi([x, y]) - \rho(x)\psi(y) = 0. \tag{4}$$

⁵ That is $v \in \mathbb{F}^n$ such that for any $B \in \mathcal{B}$, $Bv = 0$.

438 We defer the proof of Theorem 18 in Section 5.1. We now derive a corollary of Theorem 18
 439 and use it to prove Theorem 5.

440 ► **Corollary 19.** *Let $\mathfrak{g}, (\rho, V)$ satisfy the condition in Theorem 18. Then, $(\rho, V) \cong \oplus_i (\text{ad}, \mathfrak{g}_i) \oplus$
 441 (ρ', V') where \mathfrak{g}_i are not necessarily disjoint or distinct quotient algebras of \mathfrak{g} , and (ρ', V')
 442 has no linear kernel vectors.*

443 **Proof.** Let ψ be a linear kernel vector (ρ, V) . Let $V_1 = \text{im}(\psi) \cong \mathfrak{g}/\ker \psi =: \mathfrak{g}_1$. Then
 444 V_1 is invariant under ρ as for any $x \in \mathfrak{g}, \psi(y) \in V_1, \rho(x)\psi(y) = \psi([x, y]) \in V_1$. Therefore,
 445 $(\rho, V) \cong (\text{ad}, \mathfrak{g}_1) \oplus (\rho', V')$ by the semisimplicity of \mathfrak{g} . We can then repeat till we no longer
 446 have linear kernel vectors. ◀

447 ► **Theorem 5.** *Let \mathcal{B} be the image of a faithful irreducible representation ϕ of a semisimple
 448 Lie algebra \mathfrak{g} over algebraically closed fields of characteristic not 2 or 3. Then \mathcal{B} admits
 449 a linear kernel vector if and only if \mathcal{B} is trivial, or \mathfrak{g} is simple and ϕ is isomorphic to the
 450 adjoint representation.*

451 **Proof.** When \mathcal{B} is trivial it clearly has a linear vector kernel. So assume it is not. Applying
 452 Corollary 19 in the case of ρ being irreducible, we get $(\rho, V) \cong (\text{ad}, \mathfrak{g}')$ for some quotient
 453 algebra \mathfrak{g}' of \mathfrak{g} . Since ρ is faithful, we must have $\mathfrak{g} = \mathfrak{g}_i$. By definition, subrepresentations
 454 of the adjoint representation are the same as ideals. Thus irreducibility of the adjoint
 455 representation implies that \mathfrak{g} is simple. ◀

456 ► **Remark 20.** 1. Theorem 18 and Theorem 5 also hold over sufficiently large perfect fields as
 457 a semisimple Lie algebra over a perfect field remains semisimple over the algebraic closure
 458 of such fields. However, passing over to the closure need not preserve semisimplicity in
 459 general and thus, the current proof of Theorem 24 does not work for any sufficiently large
 460 field.

461 2. Initially we proved a fact equivalent to (Equation (4)) for irreducible representations
 462 of classical Lie algebras using Chevalley bases. In an attempt to simplify the proof by
 463 reducing to certain subalgebras and taking trivial subrepresentations into account, we
 464 discovered relevance of equalities (5) and (6) below. We include the previous proof in
 465 Appendix E, as some ideas and techniques there may be useful for future references.

466 5.1 Proof of Theorem 18

467 To prove Theorem 18, we need the following preparations.

468 ► **Proposition 21.** *Let $\psi : \mathfrak{g} \rightarrow V$ be a linear kernel vector for the representation (ρ, V) of a
 469 Lie algebra \mathfrak{g} . Then,*

$$470 \quad \rho(x)(\psi([x, y]) - \rho(x)\psi(y)) = 0 \tag{5}$$

471 and

$$472 \quad \rho(y)(\psi([x, y]) - \rho(x)\psi(y)) = 0 \tag{6}$$

473 *Consequently, the difference $\psi([x, y]) - \rho(x)\psi(y)$ is annihilated by $\rho(z)$ for every element z
 474 of the Lie subalgebra of \mathfrak{g} generated by x and y .*

Proof.

$$\begin{aligned}
475 \quad \rho(x)\psi([x, y]) &= -\rho([x, y])\psi(x) && \text{Applying (3) to } [x, y], x \\
476 \quad &= -\rho(x)\rho(y)\psi(x) + \rho(y)\rho(x)\psi(x) && (\rho, V) \text{ is a representation} \\
477 \quad &= -\rho(x)\rho(y)\psi(x) && \text{Applying (2) to } x \\
478 \quad &= \rho(x)\rho(x)\psi(y) && \text{Applying (3) to } x, y \\
479
\end{aligned}$$

480 Similarly (Equation (6)) follows because

$$481 \quad \rho(y)\psi([x, y]) = -\rho([x, y])\psi(y) = -\rho(x)\rho(y)\psi(y) + \rho(y)\rho(x)\psi(y) = \rho(y)\rho(x)\psi(y). \quad \blacktriangleleft$$

482 The following statement follows from standard density arguments but we provide a proof
483 in Appendix D.

484 **► Proposition 22** (Proposition 31). *Let \mathfrak{g} be an m -dimensional Lie algebra over a large
485 enough field \mathbb{F} , such that there exist two elements $x_0, y_0 \in \mathbb{F}$ that generate \mathfrak{g} . Then there
486 are elements $x_i, y_i \in \mathfrak{g}$, ($i = 1, \dots, m^2$) such that x_i and y_i generate \mathfrak{g} for every i and that
487 $x_i \otimes y_i$ span $\mathfrak{g} \otimes \mathfrak{g}$.*

488 The following lemma is a major step to prove Theorem 18.

489 **► Lemma 23.** *Let $\psi : \mathfrak{g} \rightarrow V$ be a linear kernel vector for the representation (ρ, V) of a Lie
490 algebra \mathfrak{g} over a large enough field \mathbb{F} . Suppose \mathfrak{g} can be generated by two elements, and the
491 trivial representation is not a subrepresentation of the representation (ρ, V) of \mathfrak{g} . Then for
492 every $x, y \in \mathfrak{g}$,*

$$493 \quad \psi([x, y]) - \rho(x)\psi(y) = 0.$$

494 **Proof.** By standard arguments, it is sufficient to prove the theorem for the special case
495 when \mathbb{F} is algebraically closed. We assume that. By Proposition 22, we have $\{(x_i, y_i)\}$
496 that each generate \mathfrak{g} as a Lie algebra and collectively linearly span $\mathfrak{g} \otimes \mathfrak{g}$. For each i , by
497 Proposition 21, $\rho(z)(\psi([x_i, y_i]) - \rho(x_i)\psi(y_i)) = 0$ for every $z \in \mathfrak{g}$. This equality is trilinear and
498 therefore it holds for every $z \otimes x \otimes y$ for $(z, x, y) \in \mathfrak{g} \times (\text{span}_i\{(x_i \otimes y_i)\})$. By Proposition 22,
499 $\text{span}_i\{(x_i \otimes y_i)\} = \mathfrak{g} \otimes \mathfrak{g}$ and thus, $\rho(z)(\psi([x, y]) - \rho(x)\psi(y))$ is identically zero on $\mathfrak{g}^{\otimes 3}$. Now
500 for every fixed (x, y) the vector $\psi([x, y]) - \rho(x)\psi(y)$ is annihilated by all of $\rho(\mathfrak{g})$. It follows
501 that the vector must be zero, as otherwise it would span a trivial subrepresentation. \blacktriangleleft

502 To deduce Theorem 18 from Lemma 23, we need a result for the number of generators of
503 certain Lie algebras. We recall some classical results on this topic. First, Kuranishi gave a
504 simple proof that over characteristic 0, there exists two elements that generate any semisimple
505 Lie algebra [19, Thm. 6]. The proof of the statement works directly in positive characteristic
506 (> 3) for sum of *classical simple* lie algebras, i.e. those obtained from a Chevalley basis. This
507 was extended in [2] to other kinds of simple Lie algebras over positive characteristic (> 3).
508 This can be extended to the semisimple case based on a density argument which is standard.
509 However, we couldn't find a reference for this, so we include a proof in Appendix D (see
510 Lemma 32).

511 **► Theorem 24** ([19, 2]+ Lemma 32). *Let \mathfrak{g} be a semisimple Lie algebra \mathfrak{g} over an algebraically
512 closed field \mathbb{F} of characteristic not 2 or 3. Then \mathfrak{g} can be generated by two elements.*

513 Theorem 18 follows from Lemma 23 and Theorem 24. \blacktriangleleft

514 **A** Basic notions for Lie algebras and its representations

515 A *Lie algebra* is vector space \mathfrak{g} with an alternating bilinear map, called a Lie bracket, $[-, -] :$
 516 $\mathfrak{g} \times \mathfrak{g} \rightarrow \mathfrak{g}$ that satisfies the Jacobi identity $\forall x, y, z \in \mathfrak{g}, [x, [y, z]] + [z, [x, y]] + [y, [z, x]] = 0$.
 517 A subalgebra of a Lie algebra \mathfrak{g} is a vector subspace which is closed under the Lie bracket.

518 Given two Lie algebras \mathfrak{g} and \mathfrak{h} , a *Lie algebra homomorphism* is a linear map respecting
 519 the Lie bracket, i.e., a linear map $\phi : \mathfrak{g} \rightarrow \mathfrak{h}$ such that $\phi([a, b]) = [\phi(a), \phi(b)]$.

520 Given a vector space V , we use $\mathfrak{gl}(V)$ to denote the Lie algebra, which consists of linear
 521 endomorphisms of V with the Lie bracket $[A, B] = AB - BA$ for $A, B \in \mathfrak{gl}(V)$.

522 A *representation of a Lie algebra* \mathfrak{g} is a Lie algebra homomorphism $\phi : \mathfrak{g} \rightarrow \mathfrak{gl}(V)$ for
 523 some vector space V . A subspace $U \leq V$ is *invariant* under ϕ , if for any $a \in \mathfrak{g}$, $\phi(a)(U) = U$.
 524 We say that ϕ is *irreducible*, if the only invariant subspaces under ϕ are the zero space and
 525 the full space. We say that ϕ is *completely reducible*, if there exists a proper direct sum
 526 decomposition of $V = V_1 \oplus \cdots \oplus V_c$, such that each V_i is invariant under ϕ .

527 A representation $\phi : \mathfrak{g} \rightarrow \mathfrak{gl}(V)$ is trivial, if $\phi(x) = 0 \in \mathfrak{gl}(V)$ for any $x \in \mathfrak{g}$. In this
 528 case, when V is of dimension 1, ϕ is the *trivial irreducible representation*. The adjoint
 529 representation of \mathfrak{g} , $\text{ad} : \mathfrak{g} \rightarrow \mathfrak{gl}(\mathfrak{g})$ is defined as $\text{ad}_x(y) = [x, y]$ for $x, y \in \mathfrak{g}$.

530 Suppose V is of dimension n over a field \mathbb{F} . After fixing a basis of V , $\mathfrak{gl}(V)$ can be
 531 identified as $M(n, \mathbb{F})$. Then the image of a Lie algebra representation ϕ is a matrix subspace
 532 $\mathcal{B} \leq M(n, \mathbb{F})$ that is closed under the natural Lie bracket $[A, B] = AB - BA$ for $A, B \in \mathcal{B}$.

533 **B** Correspondences between Lie algebras and Lie groups

534 Lie algebras are closely related to Lie groups. In the case of finite dimensional complex and
 535 real Lie algebras, there is a tight correspondence. Since matrix Lie algebras are the main
 536 object of study in this article, we only need results for matrix Lie algebras and matrix Lie
 537 groups, and not the most general definitions. In the following, we present some basic facts
 538 about the correspondence between Lie algebras and Lie groups in the matrix setting.

539 We follow [12] for the definitions and some basic results about matrix Lie groups and Lie
 540 algebras over \mathbb{C} that we will use later.

541 A *matrix Lie group* is a subgroup G of $GL(n, \mathbb{C})$ with the property that if $(A_m)_{m \in \mathbb{N}}$ is
 542 any sequence of matrices in G , and A_m converges to some matrix A , then either A is in G or
 543 A is non-invertible.

544 For $X \in M(n, \mathbb{C})$, define the exponential by the usual power series, that is, $e^X = \sum_{i=0}^{\infty} \frac{X^i}{i!}$.
 545 By [12, Proposition 2.1], this power series converges absolutely for any $X \in M(n, \mathbb{C})$, and e^X
 546 is a continuous function of X . A straightforward consequence of the absolute convergence is
 547 that we can differentiate term by term, which implies that $\frac{d}{dt} e^{tX} = X e^{tX} = e^{tX} X$.

548 Given a matrix Lie group G , the associated Lie algebra $\text{Lie}(G)$ is defined as $\text{Lie}(G) =$
 549 $\{X \in M(n, \mathbb{C}) \mid \forall t \in \mathbb{R}, e^{tX} \in G\}$. Let \mathfrak{g} denote $\text{Lie}(G)$; this notation is consistent with
 550 our previous notation. Clearly, for any $M \in \mathfrak{g}$, the one-parameter group $\{e^{tM} \mid t \in \mathbb{R}\}$ is a
 551 subgroup of G .

552 We need the following two classical results relating matrix Lie groups and matrix Lie
 553 algebras in Section 3.

554 **► Theorem 25** ([12, Theorem 5.20]). *Let G be a matrix Lie group with Lie algebra \mathfrak{g} and let*
 555 *\mathfrak{h} be a Lie subalgebra of \mathfrak{g} . Then there exists a unique connected Lie subgroup H of G with*
 556 *Lie algebra \mathfrak{h} . In particular, every matrix Lie algebra \mathfrak{g} is the Lie algebra of a Lie group.*

557 **► Theorem 26** ([12, Theorem 3.20 (1)]). *Let G be a matrix Lie group, and let $\mathfrak{g} = \text{Lie}(G)$.*
 558 *Then for any $X \in \mathfrak{g}$ and $g \in G$, we have $gXg^{-1} \in \mathfrak{g}$.*

C

 Some results about Cartan subalgebras

Cartan subalgebras.

Let \mathfrak{g} be a Lie algebra. A *subalgebra* $\mathfrak{h} \subseteq \mathfrak{g}$ is a vector subspace that is closed under the Lie bracket (inherited from \mathfrak{g}). In other words, $[\mathfrak{h}, \mathfrak{h}] \subseteq \mathfrak{h}$. An *ideal* $\mathfrak{i} \subseteq \mathfrak{g}$ is a subalgebra such that $[\mathfrak{g}, \mathfrak{i}] \subseteq \mathfrak{i}$. Let \mathfrak{g}_1 and \mathfrak{g}_2 be ideals of \mathfrak{g} . Define $[\mathfrak{g}_1, \mathfrak{g}_2] = \text{span}([x, y] \mid x \in \mathfrak{g}_1, y \in \mathfrak{g}_2)$. Let $\mathfrak{g}^1 = \mathfrak{g}$ and inductively define $\mathfrak{g}^i = [\mathfrak{g}^{i-1}, \mathfrak{g}]$. An algebra \mathfrak{g} is called *nilpotent* if there is an n such that $\mathfrak{g}^n = 0$. Similarly, define $\mathfrak{g}^{(1)} = \mathfrak{g}$ and $\mathfrak{g}^{(i)} = [\mathfrak{g}^{(i-1)}, \mathfrak{g}^{(i-1)}]$. An algebra \mathfrak{g} is called *solvable* if there is an n such that $\mathfrak{g}^{(n)} = 0$. The *normalizer* of a subspace \mathfrak{a} of \mathfrak{g} is defined as $\mathfrak{n}_{\mathfrak{g}}(\mathfrak{a}) = \{x \in \mathfrak{g} \mid [x, \mathfrak{a}] \subseteq \mathfrak{a}\}$. A subalgebra \mathfrak{h} of \mathfrak{g} is a *Cartan subalgebra* if it is nilpotent and $\mathfrak{n}_{\mathfrak{g}}(\mathfrak{h}) = \mathfrak{h}$.

We shall need the following classical result on Cartan subalgebras. For $x \in \mathfrak{g}$, recall that $\text{ad}_x : \mathfrak{g} \rightarrow \mathfrak{g}$ is the linear map defined by $\text{ad}_x(y) = [x, y]$ for $y \in \mathfrak{g}$. In particular, the exponentiation e^{ad_x} is a linear map from \mathfrak{g} to \mathfrak{g} , and it is a Lie algebra automorphism if ad_x is nilpotent, called an inner automorphism. The group generated by inner automorphisms is denoted by $\text{Int}(\mathfrak{g})$.

► **Theorem 27** (See e.g. [4, Chapter 3.5]). *Let \mathfrak{g} be a Lie algebra over an algebraically closed field \mathbb{F} of characteristic zero. For any two Cartan subalgebras \mathfrak{h}_1 and \mathfrak{h}_2 , there exists $g \in \text{Int}(\mathfrak{g})$ such that $\mathfrak{h}_1 = g(\mathfrak{h}_2)$.*

To recover the statement in Lemma 14, note that for a matrix Lie algebra $\mathcal{B} \leq M(n, \mathbb{C})$, an inner automorphism takes the form as a conjugation by an invertible matrix. This is because $\text{Ad}(e^x) = e^{\text{ad}_x}$, where Ad is the conjugation by matrices. This can be seen, e.g., by taking the derivative of $\text{Ad}(e^x)Y = e^{tx}Y e^{-tx}$ at $t = 0$.

► **Theorem 28** (Lie's theorem on solvable Lie algebras). *Let \mathbb{F} be an algebraically closed field of characteristic zero. Let $\mathcal{B} \leq M(n, \mathbb{F})$ be a solvable matrix Lie algebra over \mathbb{F} . Then there exists $T \in \text{GL}(n, \mathbb{F})$, such that for any $B \in \mathcal{B}$, TBT^{-1} is upper triangular.*

Regular elements of Lie algebras.

Let λ be a formal variable, and let $\sum_i c_{i,x} \lambda^i$ be the characteristic polynomial of ad_x . The smallest r such that $c_{r,x}$ is not identically zero over all $x \in \mathfrak{g}$ is called the *rank* of \mathfrak{g} . The open set of points $\{x \in \mathfrak{g} \mid c_r(x) \neq 0\}$ is the set of *regular points*. A simple observation is that the set of regular elements is Zariski open and thus it is dense.

For $x \in \mathfrak{g}$, the Fitting null component of ad_x is $F_0(\text{ad}_x) = \{y \in \mathfrak{g} : \exists m > 0, \text{ad}_x^m(y) = 0\}$.

Regular elements and Cartan subalgebras are closely related as the following theorem shows.

► **Theorem 29** ([4, Corollary 3.2.8]). *Let \mathfrak{g} be a Lie algebra over a field of order larger than $\dim(\mathfrak{g})$. For a regular $x \in \mathfrak{g}$, $F_0(\text{ad}_x)$ is a Cartan subalgebra.*

Computing Cartan subalgebras.

We shall need the following result of de Graaf, Ivanyos, and Rónyai [5] regarding computing Cartan subalgebras. In algorithms, Lie algebras are often given by structure constants. That is, let \mathfrak{g} be a Lie algebra of dimension n over a field \mathbb{F} , and let a_1, \dots, a_n be a linear basis of \mathfrak{g} . The *structure constants* α_{ijk} ($i, j, k \in \{1, \dots, n\}$) are field elements such that $[a_i, a_j] = \sum_{k=1}^n \alpha_{ijk} a_k$.

600 ▶ **Theorem 30** ([5, Theorem 5.8]). Let \mathfrak{g} be a Lie algebra of dimension n over a field \mathbb{F} with
 601 $|\mathbb{F}| > n$. Suppose \mathfrak{g} is given by its structure constants with respect to a basis a_1, \dots, a_n , and fix
 602 $\Omega \subseteq \mathbb{F}$ such that $|\Omega| = n + 1$. Then there is a deterministic polynomial-time algorithm which
 603 computes a regular element $x = \sum \alpha_i a_i$, $\alpha_i \in \Omega$, such that $F_0(\text{ad}_x)$ is a Cartan subalgebra of
 604 \mathfrak{g} .

605 Note that to obtain Theorem 13, we start with a matrix Lie algebra $\mathcal{B} = \langle B_1, \dots, B_m \rangle \leq$
 606 $M(n, \mathbb{F})$, compute structure constants by expanding $[B_i, B_j] = \sum_{k \in [m]} \alpha_{i,j,k} B_k$, apply The-
 607 orem 30, and use its output to obtain a subspace of \mathcal{B} which is a Cartan subalgebra.

608 **D** Density arguments and the generation of Lie algebras

609 In this part we prove some facts based on standard density arguments.

610 Let U be an m -dimensional vector space over an infinite field \mathbb{F} . By choosing a basis
 611 we identify U with \mathbb{F}^m . We say that a nonempty subset D of U is *huge* if there exists a
 612 nonzero polynomial $f(t_1, \dots, t_m) \in \mathbb{F}[t_1, \dots, t_m]$ such that if $u = (u_1, \dots, u_m)^T \notin D$ then
 613 $f(u_1, \dots, u_m) = 0$. (Thus, huge subsets are those that contain Zariski open subsets.) It is
 614 easy to see that hugeness is independent of the choice of the basis and that the intersection
 615 of finitely many huge subsets is huge as well. As a hyperplane of U consists of the zeros of a
 616 linear function on U , we have that any huge subset of U spans U .

617 Let \mathfrak{g} be an m -dimensional Lie algebra over \mathbb{F} . Let u_1, \dots, u_m be a basis for \mathfrak{g} . Recall
 618 that the structure constants α_{ijk} ($i, j, k \in \{1, \dots, m\}$) are field elements such that such
 619 that $[u_i, u_j] = \sum_{k=1}^m \alpha_{ijk} u_k$. A *Lie expression* or Lie polynomial $E(z_1, \dots, z_\ell)$ in ℓ variables
 620 z_1, \dots, z_ℓ is an expression that can be recursively built using linear combinations and the
 621 bracket symbol. Let $x_i = \sum_{j=1}^m x_{ij} u_j$ ($i = 1, \dots, \ell$). Then the structure constants can
 622 be used to expand $E(x_1, \dots, x_\ell)$ as a vector whose coordinates are polynomials in x_{ij} . If
 623 we assign $m - 1$ elements of \mathfrak{g} to the variables z_2, \dots, z_ℓ then E expands to an a vector
 624 whose coordinates are polynomials in $x_{11}, \dots, x_{1\ell}$ that may include nonzero constant terms.
 625 Therefore it will be convenient to also consider Lie expressions *over* \mathfrak{g} : these are expressions
 626 which may include constant elements from \mathfrak{g} . From the definition of density it follows
 627 that if E is an expression in a single variable z that is not identically zero on \mathfrak{g} then the
 628 elements x of \mathfrak{g} on which E evaluates to a nonzero element of \mathfrak{g} is huge. Furthermore, if
 629 there are m expressions $E_1(z), \dots, E_m(z)$ such that there exists an element $x \in \mathfrak{g}$ such that
 630 $E_1(x), \dots, E_m(x)$ are linearly independent then such elements are a huge subset of \mathfrak{g} . To see
 631 this, just consider the determinant expressing that $E_1(x) \dots, E_m(x)$ are linearly dependent.

632 ▶ **Proposition 31.** Let \mathfrak{g} be an m -dimensional Lie algebra over a large enough field \mathbb{F} such
 633 that there exist two elements $x_0, y_0 \in \mathbb{F}$ that generate \mathfrak{g} . Then there are elements $x_i, y_i \in \mathfrak{g}$,
 634 ($i = 1, \dots, m^2$) such that x_i and y_i generate \mathfrak{g} for every i and that $x_i \otimes y_i$ span $\mathfrak{g} \otimes \mathfrak{g}$.

635 **Proof.** Pick expressions $E_i(z, w)$ ($i = 1, \dots, m$) such that $E_i(x_0, y_0)$ are linearly independent.
 636 Then the set of elements x such that $E_i(x, y_0)$ are linearly independent is huge and hence
 637 contains a basis x_1, \dots, x_m of \mathfrak{g} . For each j , the subset consisting of those y for which
 638 $E_i(x_j, y)$ are linearly independent is a huge set $D_j \subset \mathfrak{g}$ whence there exist elements $y_{jk} \in D_j$
 639 ($k = 1, \dots, m$) that are a basis for \mathfrak{g} . Each of the m^2 pairs x_j, y_{jk} generate \mathfrak{g} . To see that
 640 they span $\mathfrak{g} \otimes \mathfrak{g}$, write an element $z \in \mathfrak{g} \otimes \mathfrak{g}$ in the form $z = \sum x_j \otimes y'_j$ and express y'_j as
 641 $y'_j = \sum_k \alpha_{jk} y_{jk}$. Then $z = \sum_{j,k} \alpha_{jk} x_j \otimes y_{jk}$. ◀

642 ▶ **Lemma 32.** Let $\mathfrak{g}_1, \dots, \mathfrak{g}_m$ be finite dimensional simple Lie algebras over a large enough
 643 field \mathbb{F} , each generated by 2 elements. Then $\mathfrak{g}_1 \oplus \dots \oplus \mathfrak{g}_m$ is also generated by two elements.

644 **Proof.** Assume that x_i and y_i generate \mathfrak{g}_i ($i = 1, \dots, m$). We claim that we may further
 645 assume that $\text{ad}(x)^{d_i} y_i \neq 0$ where $d_i = \dim_{\mathbb{F}}(\mathfrak{g}_i)$. Indeed, if $\text{ad}(x)^{d_i} y_i = 0$ then, by Engel's
 646 theorem, there exists a pair $(w_i, z_i) \in \mathfrak{g} \times \mathfrak{g}$ such that $\text{ad}(w_i)^{d_i} z_i \neq 0$. If we fix z_i from
 647 such a pair then the elements w_i such that $\text{ad}(w_i)^{d_i} z_i \neq 0$ is a huge set. There exist d_i Lie
 648 expressions E_1, \dots, E_{d_i} in in two variables such that $E_1(x_i, y_i), \dots, E_{d_i}(x_i, y_i)$ are linearly
 649 independent elements of \mathfrak{g}_i . The elements w_i such that $E_1(w_i, y_i), \dots, E_{d_i}(w_i, y_i)$ are linearly
 650 independent form a huge set. The intersection of these two huge sets is still huge and hence
 651 non-empty. We replace x_i with an element from the intersection. Now the set of w_i such
 652 that $\text{ad}(x_i)^{d_i} w_i \neq 0$ is huge as well as these set of those for which $E_j(x_i, w_i)$ are linearly
 653 independent. We can replace y_i with an element from the intersection.

654 Let $f_i = f_i(t)$ be the monic polynomial of smallest degree such that $f_i(\text{ad}(x_i))y_i = 0$.
 655 Note that f_i has degree at most d_i and the assumption on x_i and y_i implies that f_i is not a
 656 divisor of t^{d_i} . Therefore each f_i has a nonzero root (in the algebraic closure $\overline{\mathbb{F}}$ of \mathbb{F}). Let
 657 R_i be the set of nonzero roots of f_i in $\overline{\mathbb{F}}$. There exist field elements $\alpha_1, \dots, \alpha_m \in \mathbb{F}$ such
 658 that the sets $\alpha_i R_i$ are pairwise disjoint. Replacing x_i with $\alpha_i x_i$ we arrange that the sets R_i
 659 become pairwise disjoint. Then for each i , put $h_i = \prod_{j \neq i} f_j$. We have that $h_i(\text{ad}(x_j))y_j = 0$
 660 for every $j \neq i$, while $h_i(\text{ad}(x_j))y_i \neq 0$ as h_i is not divisible by f_i .

661 Put $x = \sum_{i=1}^m x_i$ and $y = \sum_{i=1}^m y_i$. Then $h_i(x)y$ is a nonzero element of \mathfrak{g}_i . Let M
 662 be the subalgebra of \mathfrak{g} generated by x and y . We see that M has a nonzero element, say
 663 z_i contained in \mathfrak{g}_i . The projection of M on the i th component is clearly \mathfrak{g}_i and x_i and y_i
 664 generate \mathfrak{g}_i . As \mathfrak{g}_i is simple we have that the ideal of M generated by z_i is \mathfrak{g}_i . This holds for
 665 all $i = 1, \dots, m$, showing that $M = \mathfrak{g}$. ◀

666 **E** Linear kernel vectors of matrix Lie algebras

667 In this section, we will give an alternative proof of Lemma 23, which doesn't use density
 668 arguments, but instead uses weight decomposition of representations of semi-simple Lie
 669 algebras over \mathbb{C} .

670 **E.1** Weight decomposition of Lie algebra representations

671 Fix a Cartan subalgebra \mathfrak{h} of a semisimple Lie algebra \mathfrak{g} over \mathbb{C} . By definition \mathfrak{h} is nilpotent.
 672 If \mathfrak{g} is semisimple, \mathfrak{h} is abelian [24, Thm3, Ch.3]. Similar to the notion of eigenvalues and
 673 eigenspaces is the concept of a weight and its weight space. Intuitively, it can be thought
 674 of as a linear function that captures the eigenvalues of a set of matrices simultaneously.
 675 Formally, a *weight* is an element of \mathfrak{h}^* . If $w \in \mathfrak{h}^*$, then the w -weight space of V is defined as
 676 $V_w = \{v \in V \mid \forall h \in \mathfrak{h}, \rho(h) \cdot v = w(h)v\}$.

677 ▶ **Theorem 33.** *If \mathfrak{g} is a complex semi-simple Lie algebra, then every representation (ρ, V)
 678 can be decomposed into weight spaces $V = \bigoplus_w V_w$.*

679 Using this decomposition we have a basis such that for any $h \in \mathfrak{h}$, $\rho(h)$ is a diagonal
 680 matrix with $w(h)$ as diagonal elements where w runs over all weights of V .

681 ▶ **Fact 34.** *The matrix space defined by the image $\rho(\mathfrak{g})$ of the representation (ρ, V) is singular
 682 iff 0 is a weight of the representation, i.e. V_0 occurs with multiplicity at least one.*

683 **Proof.** This follows easily from the observation about $\rho(h_i)$ which implies that $\rho(\mathfrak{h})$ is singular
 684 if 0 is a weight. From Section 4, we know that the entire algebra is singular if any of its
 685 Cartan subalgebra is. If 0 is not a weight, then it is easy to construct a element $h \in \mathfrak{h}$ such
 686 that $w(h) \neq 0$ for every weight. Thus, $\rho(h)$ has full rank. ◀

687 If we decompose the adjoint representation, the weights we obtain are called *roots* usually
 688 denoted by Φ . It is also a fact that if $\alpha \in \Phi$, then $-\alpha \in \Phi$. We thus can write $\mathfrak{g} = \mathfrak{h} \oplus_{\alpha \in \Phi} \mathfrak{g}_\alpha$.
 689 Moreover, each of the spaces \mathfrak{g}_α is one-dimensional. We denote an element of \mathfrak{g}_α by g_α which
 690 is unique upto a scalar. Such a decomposition of the Lie algebra is very useful as we can
 691 understand the action under any representation of these subspaces \mathfrak{g}_α as follows .

692 ► **Proposition 35.** [24, Prop.1, Chapter 7] For any representation (ρ, V) of \mathfrak{g} , $\rho(g_\alpha)V_w \subset$
 693 $V_{w+\alpha}$ for every weight w and every root α .

694 E.2 Notation

695 Fix a complex semi-simple Lie algebra \mathfrak{g} and a Cartan subalgebra \mathfrak{h} . Let Φ be its set of roots.
 696 We choose a Cartan-Weyl basis⁶ for \mathfrak{g} (cf. e.g. [24, pp. 48]). This means that we have a set
 697 of simple roots⁷ $S = \{\alpha_1, \dots, \alpha_n\}$ and a basis of \mathfrak{h} , $\{h_1, \dots, h_n\}$ such that the following hold,

$$\begin{aligned} 698 \quad [h_i, g_{\alpha_j}] &= \alpha_j(h_i)g_{\alpha_j} && \forall i, j \in [n] \\ 699 \quad [g_\alpha, g_\beta] &= c_{\alpha\beta}g_{\alpha+\beta} \quad (c_{\alpha\beta} \neq 0) && \alpha + \beta \in \Phi \\ 700 \quad [g_\alpha, g_\beta] &= 0 && \text{if } \alpha + \beta \notin \Phi \\ 701 \quad [g_{\alpha_i}, g_{-\alpha_i}] &= h_i && \forall i \in [n] \end{aligned}$$

703 Choose a basis of $V \cong \mathbb{F}^N$, such that $\rho(\mathfrak{h})$ is diagonal. Let W be the set of weights of V
 704 and thus $V = \bigoplus_{w \in W} V_w$ such that V_w is the w weight space of \mathfrak{h} . Note that we are assuming
 705 that $0 \in W$ as non-singular spaces anyway cannot have a linear kernel.

706 E.3 Main proof

707 We recall that given a Lie algebra \mathfrak{g} and a representation (ρ, V) a linear kernel vector
 708 $\phi : \mathfrak{g} \rightarrow V$ is a linear map such that $\rho(x)\phi(x) = 0$ for every $x \in \mathfrak{g}$. We state the main lemma
 709 we need and will prove it later.

710 ► **Lemma 36.** Assume that trivial representation is not a subrepresentation of the represent-
 711 ation (ρ, V) of \mathfrak{g} . Let $\psi : \mathfrak{g} \rightarrow V$ be a linear kernel vector. Then for any $\alpha, \beta \in \Phi$ such that
 712 $\alpha + \beta \neq 0$ and $h \in \mathfrak{h}$ we have

$$\begin{aligned} 713 \quad \psi(h) &\in V_0, \psi(g_\alpha) \in V_\alpha \text{ and} \\ 714 \quad \psi([g_\alpha, g_\beta]) &= \rho(g_\alpha)\psi(g_\beta) \\ 715 \quad \psi([h, g_\alpha]) &= \rho(h)\psi(g_\alpha) \end{aligned}$$

717 ► **Theorem 37.** Assume that trivial representation is not a subrepresentation of the repres-
 718 entation (ρ, V) of \mathfrak{g} . Then for every $x, y \in \mathfrak{g}$,

$$719 \quad \psi([x, y]) - \rho(x)\psi(y) = 0$$

720 **Proof.** By linearity, it suffices to show this for the basis vectors h_i, g_α . Lemma 36 shows
 721 it in every except when we have $(x, y) = (g_\alpha, g_{-\alpha})$. Fix any root α . By Lemma 36 and
 722 Proposition 35, the vector $\psi([g_\alpha, g_{-\alpha}]) - \rho(g_\alpha)\psi(g_{-\alpha}) \in V_0$ and thus is annihilated by $\rho(\mathfrak{h})$.

⁶ Cartan-Weyl basis and the Chevally basis differ only by a normalization. We do not need properties of the coefficients $c_{\alpha,\beta}$ and thus either basis works well.

⁷ Simple roots are just a basis of \mathfrak{h}^* while the set of all roots can be linearly dependent.

723 We now wish to show that it is annihilated by $\rho(g_\beta)$ for any root β . The assumption that
 724 there are no trivial subrepresentations then implies that this vector must be zero.

725 Proposition 21 already shows it if $\beta \in \{\alpha, -\alpha\}$ and so we assume that's not the case.

$$\begin{aligned}
 726 \quad \rho(g_\beta)\rho(g_\alpha)\psi(g_{-\alpha}) &= \rho([g_\beta, g_\alpha])\psi(g_{-\alpha}) + \rho(g_\alpha)\rho(g_\beta)\psi(g_{-\alpha}) && \rho \text{ is a representation} \\
 727 \quad &= \rho([g_\beta, g_\alpha])\psi(g_{-\alpha}) + \rho(g_\alpha)\psi([g_\beta, g_{-\alpha}]) && \text{Lemma 36 for } (\beta, -\alpha) \\
 728 \quad &= \psi([[g_\beta, g_\alpha], g_{-\alpha}]) + \psi([g_\alpha, [g_\beta, g_{-\alpha}]]) && \text{36 for } (\alpha + \beta, -\alpha), (\alpha, \beta - \alpha) \\
 729 \quad &= \psi([g_\beta, g_\alpha], g_{-\alpha}) + [g_\alpha, [g_\beta, g_{-\alpha}]] && \text{By linearity} \\
 730 \quad &= \psi([g_\beta, [g_\alpha, g_{-\alpha}]]) && \text{By Jacobi identity} \\
 731 \quad &= \rho(g_\beta)\psi([g_\alpha, g_{-\alpha}]) && \text{36 for } (\beta, h), h = [g_\alpha, g_{-\alpha}] \in \mathfrak{h}
 \end{aligned}$$

733 Here, we have used Lemma 36 formally even if one of them is not a root to prevent dividing
 734 into cases. For example, if $\beta - \alpha$ is not a root then that term is anyway 0 and we can
 735 represent 0 as $\psi(0) = \psi([g_\beta, g_{-\alpha}])$. ◀

736 E.4 Proof of Lemma 36

737 For ease of notation, we label $H_i = \rho(h_i)$ for $1 \leq i \leq n$ and $X_\alpha = \rho(g_\alpha)$, $\alpha \in \Phi$. Similarly,
 738 we will have $v_i := \psi(h_i)$, $v_\alpha := \psi(g_\alpha)$. We restate Equation (3) in more verbose terms,

$$739 \quad H_i v_i = 0 \quad \forall i \in [n] \quad (7)$$

$$740 \quad X_\alpha v_\alpha = 0 \quad \forall \alpha \in \Phi \quad (8)$$

$$741 \quad H_i v_j + H_j v_i = 0 \quad \forall i, j \in [n], i \neq j \quad (9)$$

$$742 \quad H_i v_\alpha + X_\alpha v_i = 0 \quad \forall i \in [n], \alpha \in \Phi \quad (10)$$

$$743 \quad X_\beta v_\alpha + X_\alpha v_\beta = 0 \quad \forall \alpha, \beta \in \Phi \quad (11)$$

745 ▶ **Lemma 38 (Structure).** *For every $i \in [n]$, $v_i \in V_0$, and for every root α , $v_\alpha \in V_\alpha$ if α is a
 746 weight and is 0 otherwise.*

747 **Proof.** i) Let $v_j = \sum_{w \in W} u_w$ where $u_w \in V_w$. Since $0 = H_j v_j = \sum_{w \in W} w(h_j)u_w$, we have
 748 that $w(h_j)u_w = 0$. For any $w \neq 0$ such that $u_w \neq 0$ we have $w(h_j) = 0$. Pick $k \in [n]$ such that
 749 $w(h_k) \neq 0$. Then, $H_k v_j + H_j v_k = 0$. Looking at the V_w component $w(h_k)u_w + w(h_j)v_k = 0$.
 750 Since, $w(h_j) = 0$, we get that $w(h_k)u_w = 0$. But k is chosen such that $w(h_k) \neq 0$ and thus,
 751 $u_w = 0$.

752 ii) Fix an α . We just proved that $\forall i, v_i \in V_0$ and using Proposition 35 we get $X_\alpha v_i \in V_\alpha$.
 753 Suppose $v_\alpha = \sum_{w \in W} u_w$, where $u_w \in V_w$. For any non-zero $w \neq \alpha$, pick i such that
 754 $w(h_i) \neq 0$. Then, $H_i v_\alpha + X_\alpha v_i = 0$. But we already know that $X_\alpha v_i \in V_\alpha$ and thus
 755 $H_i v_\alpha \in V_\alpha$. Suppose $H_i v_\alpha = \sum_{w \in W} w(h_i)u_w \in V_\alpha$.

756 Then, the V_w component should be zero but $w(h_i) \neq 0 \implies u_w = 0$. It follows that
 757 $v_\alpha \in V_\alpha \oplus V_0$ for every root α . Fix α and now for every $\beta \neq \alpha$, we have $X_\alpha v_\beta + X_\beta v_\alpha = 0$.
 758 Comparing the V_β component we get that $X_\beta v_\alpha = 0$. This is true for every β and since
 759 $u_0 \in V_0$, it is also true for $\rho(h)$ for every $h \in \mathfrak{h}$. Thus, for every $x \in \mathfrak{g}$, $\rho(x)u_w$ and since
 760 there are no trivial submodules, $u_0 = 0$. Thus, $v_\alpha \in V_\alpha$. ◀

761 ▶ **Lemma 39.** *For all pairs of roots α, β such that $0 \neq \beta + \alpha \in \Phi$, $X_\beta v_\alpha = c_{\alpha\beta} v_{\alpha+\beta}$.*

762 **Proof.** We have that $H_i v_{\alpha+\beta} + X_{\alpha+\beta} v_i = 0$. Similarly, $X_\alpha v_i = -H_i v_\alpha = -\alpha(h_i)v_\alpha$ and
 763 $X_\beta v_i = -H_i v_\beta = -\beta(h_i)v_\beta$. Moreover, $X_\beta v_\alpha + X_\alpha v_\beta = 0$.

764 Now, $X_{\alpha+\beta} = \rho(g_{\alpha+\beta}) = c\rho([g_\alpha, g_\beta]) = c(X_\alpha X_\beta - X_\beta X_\alpha)$ where $c = \frac{1}{c_{\alpha\beta}}$. Thus,

$$\begin{aligned}
 765 \quad & -(\alpha + \beta)(h_i)v_{\alpha+\beta} = X_{\alpha+\beta}v_i \\
 766 \quad & \quad = c(X_\alpha X_\beta - X_\beta X_\alpha)v_i \\
 767 \quad & \quad = c(X_\alpha(X_\beta v_i) - X_\beta(X_\alpha v_i)) \\
 768 \quad & \quad = c(X_\alpha(-\beta(h_i)v_\beta) - X_\beta(-\alpha(h_i)v_\alpha)) \\
 769 \quad & \quad = -c(\beta(h_i)X_\alpha v_\beta - \alpha(h_i)X_\beta v_\alpha) \\
 770 \quad & \quad = -c(-\beta(h_i)X_\beta v_\alpha - \alpha(h_i)X_\beta v_\alpha) \\
 771 \quad & \quad = c \cdot (\alpha + \beta)(h_i) \cdot X_\beta v_\alpha.
 \end{aligned}$$

773 Thus, picking h_i such that $(\alpha + \beta)(h_i) \neq 0$, we get that $X_\beta v_\alpha = \frac{1}{c}v_{\alpha+\beta}$. ◀

774 ▶ **Lemma 40.** *Let α, β be roots such that $\beta \neq -\alpha, \alpha + \beta \notin \Phi$. Then, we have that $X_\alpha v_\beta = 0$.*

775 **Proof.** Since, $\alpha + \beta \notin \Phi$, $[g_\alpha, g_\beta] = 0$ which implies that $\rho([g_\alpha, g_\beta]) = 0$ and thus, $X_\alpha X_\beta =$
 776 $X_\beta X_\alpha$. Moreover, $\exists h_i$ such that $\alpha(h_i) \neq -\beta(h_i)$ because the h_i form a basis for \mathfrak{h} . We fix our i
 777 to be one such. Now, from Equation (11), we get that $X_\alpha v_\beta + X_\beta v_\alpha = 0$. From Equation (10),
 778 we get that $H_i v_\alpha + X_\alpha v_i = 0$ and multiplying it by X_β we obtain, $\alpha(h_i)X_\beta v_\alpha + X_\beta X_\alpha v_i = 0$.
 779 Repeating it with β and α switched, we get, $\beta(h_i)X_\alpha v_\beta + X_\alpha X_\beta v_i = 0$. Subtracting these
 780 2 equations, we get $\alpha(h_i)X_\beta v_\alpha - \beta(h_i)X_\alpha v_\beta = 0$. We already have another equation i.e.
 781 $X_\alpha v_\beta + X_\beta v_\alpha = 0$. Since $\beta(h_i) \neq -\alpha(h_i)$, these two homogeneous equations are independent
 782 and thus, the only solution is that $X_\alpha v_\beta = X_\beta v_\alpha = 0$. ◀

783 The structure lemma establishes Lemma 36 when $y \in \mathfrak{h}$ i.e. for any $x \in \mathfrak{g}, h \in \mathfrak{h}$ we
 784 have $\rho(h)\psi(x) = \psi([h, x])$. To see this notice that $\rho(h)\psi(g_\alpha) = \alpha(h)\psi(g_\alpha) = \psi(\alpha(h)g_\alpha) =$
 785 $\psi([h, g_\alpha])$ where the first equality uses that $\psi(g_\alpha) \in V_\alpha$ and the last by the property of the
 786 basis. And the other two lemmas extend it to α, β as $c_{\alpha\beta}v_{\alpha+\beta} = c_{\alpha\beta}\psi(g_{\alpha+\beta}) = \psi([g_\alpha, g_\beta])$.

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